

Lecture 6III While-Looping

- Binary search
- While loop
 - Root Finding

An array of sorted integers

- X is an array of n sorted integers
- $X[i] < X[j]$, if $i < j$
- $X = [1\ 3\ 5\ 7\ 9\ 11\ 13\ 15\ 17]$
- e.x. $X[3] = 5 < X[6] = 11$

Searching

- s is an integer
- Is s in X ?
- Sequential search: Compare s with $X[1], X[2], \dots, X[n]$ one by one
- Time consuming

- For large n , it is time consuming to compare s with elements in X one by one
- Binary search improves this drawback
- Binary search is more efficient than sequential search

An array of sorted elements

- X is a vector of positive integers
- $X[i] < X[j]$ if $i < j$
- Input s
- Output i
 - $i > 0$, where $x[i] == s$
 - $i=0$, if $s \neq x[i]$ for all i

Example

- $X=[1\ 3\ 5\ 7\ 9\ 11\ 13\ 15\ 17]$ and $s=5$
- Output : 3

- $X=[1\ 3\ 5\ 7\ 9\ 11\ 13\ 15\ 17]$ and $s=4$
- Output : 0

Sorting

```
N=10;  
X=ceil(rand(1,N)*10000);  
X=sort(X)
```

X =

Columns 1 through 7

1704	1834	2141	2633	6022	6050	6366
------	------	------	------	------	------	------

Columns 8 through 10

6596	6597	7537
------	------	------

```
N=500;  
X=ceil(rand(1,N)*10000);  
X=sort(X)
```

```
a=1;b=length(X);
```


- Binary search cuts $[a, b]$ into two sub-interval

$$[a, c - 1] \cup \{c\} \cup [c + 1, b]$$

- Where $c = \text{floor}((a+b)/2)$

- Halting condition: $s == X[c] \mid a > b$
- If $s < X[c]$, the answer should belong $[a, c-1]$

$$b \leftarrow c - 1$$

- If $X[c] < s$, the answer should belong $[c+1, b]$

$$a \leftarrow c + 1$$

left interval

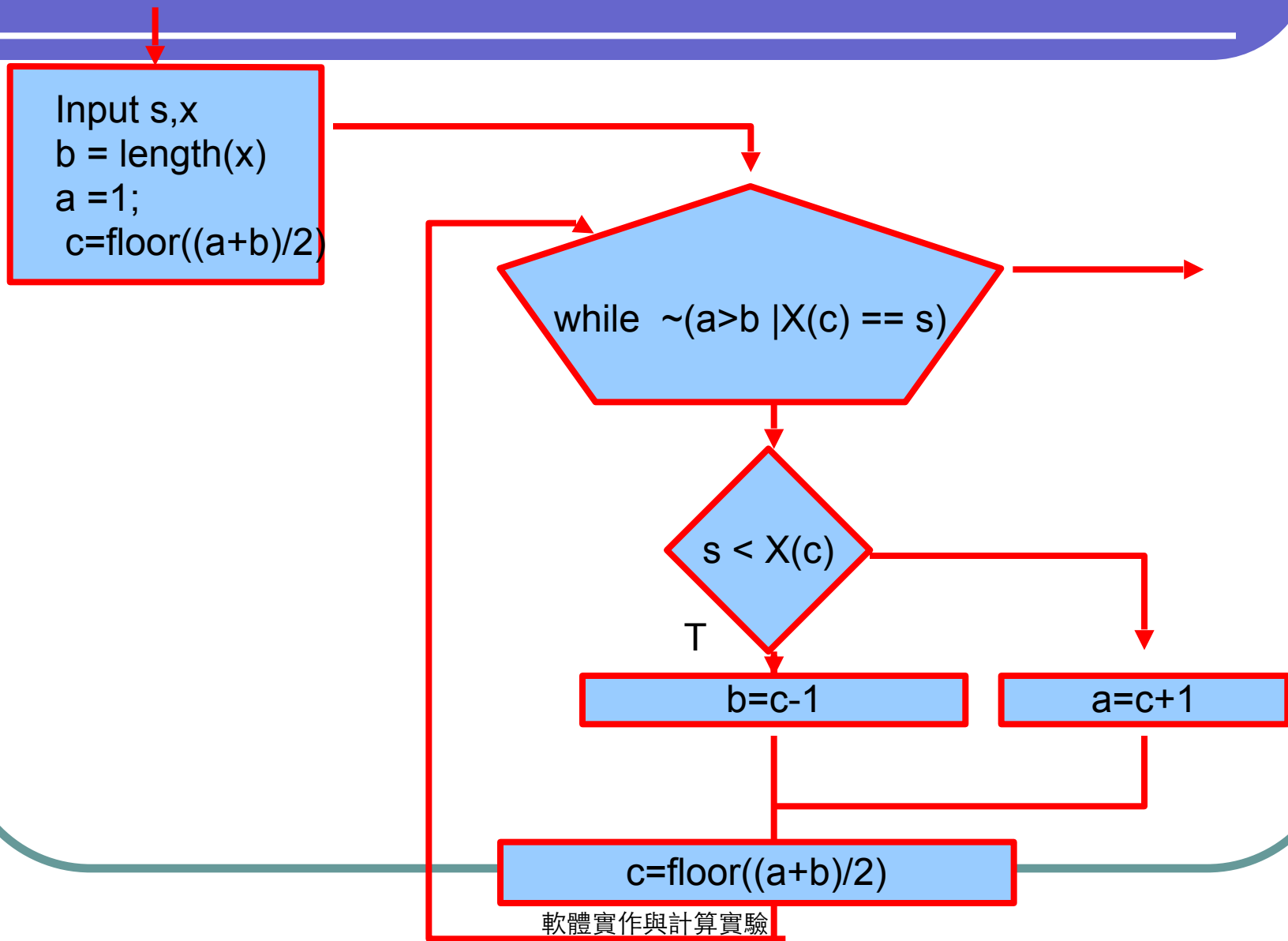
- $X=[1\ 3\ 5\ 7\ 9\ 11\ 13\ 15\ 17]$ and $s=4$
- $a=1; b=9$
- $c=5$
- The location of s is within $[a,c]$
- By operation, $b \leftarrow c - 1$
the searching interval $[a\ b]$ becomes $[1\ 4]$

Right interval

- $X=[1\ 3\ 5\ 7\ 9\ 11\ 13\ 15\ 17]$ and $s=13$
- $a=1; b=9$
- $c=5$
- The location of s is within $[c,b]$
- By operation $a \leftarrow c + 1$
the searching interval $[a\ b]$ becomes $[6\ 9]$

Right interval

- $X=[1\ 23\ 45\ 38\ 66\ 77\ 88\ 99\ 101\ 999]$ and $s=66$
- $a=1; b=5$
- $c=3$
- The location of s is within $[c,b]$
- By operation $a \leftarrow c$
the searching interval $[a\ b]$ becomes $[3\ 5]$



- Entry condition

$$\sim(a > b \mid X(c) == s)$$

- Halting condition

$$(a > b \mid X(c) == s)$$

Inline : $f(x) = \sin(x)$

```
fs = input('f(x) = ','s');  
f = inline(fs);
```

```
>> f(pi/2)
```

```
ans =
```

```
1
```


Roots

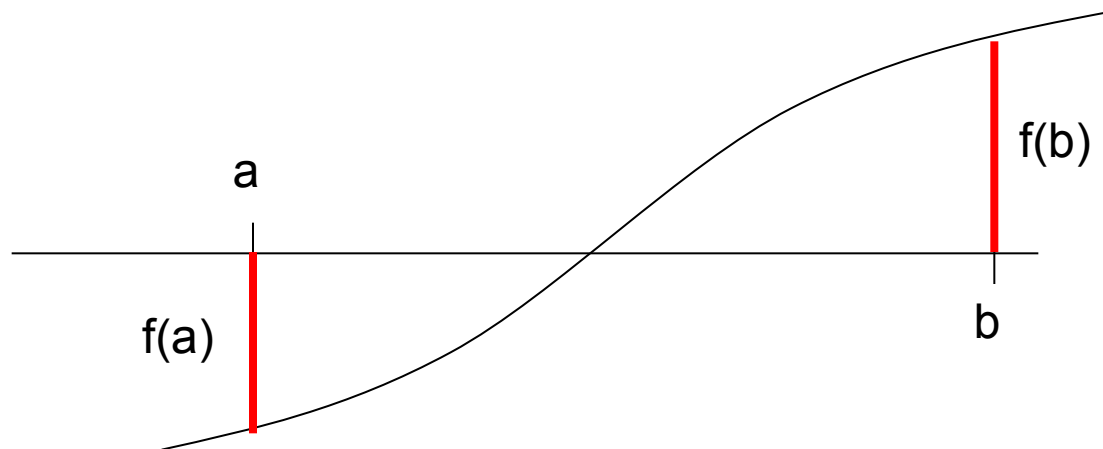
- Mathematics
 - x is a root of $f(x)$ if $f(x) = 0$
- Numerical analysis
 - x is a root of $f(x)$ if $\text{abs}(f(x)) < \text{eps}$

```
>> sin(pi)  
  
ans =  
  
1.2246e-016
```

```
>> abs(sin(pi)) < eps  
  
ans =  
  
1
```

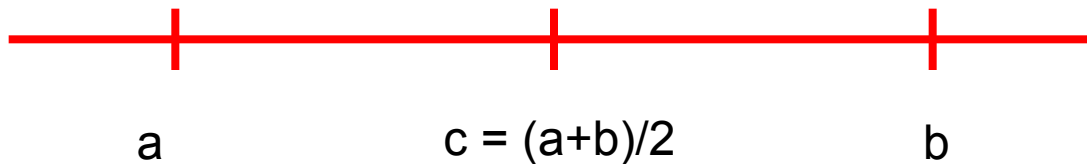
Existence of roots

- $f(x)$ is well defined over interval $[a,b]$
- If $f(x)$ is continuous and $f(a)f(b) < 0$ there exists at least one root within $[a,b]$



Bipartition

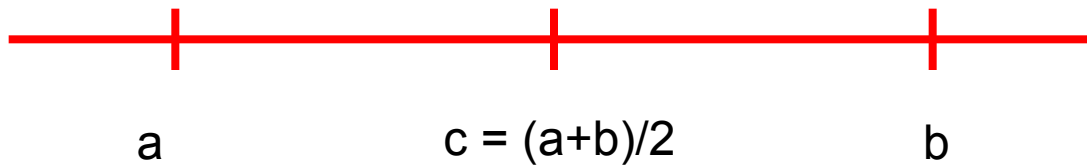
- Partition interval $[a,b]$ to two intervals such that $[a,b]=[a,c] \cup [c,b]$, where $c = (a+b) / 2$



Proposition

If $f(a)f(b) < 0$

it holds that $f(a)f(c) < 0$ or $f(c)f(b) < 0$



Proof

$$f(a)f(b) < 0, f(c) \neq 0$$

$$(I) f(a) > 0 \text{ and } f(b) < 0$$

$$\Rightarrow f(c)f(a) < 0 \text{ or } f(c)f(b) < 0$$

$$(II) f(a) < 0 \text{ and } f(b) > 0$$

$$\Rightarrow f(c)f(a) < 0 \text{ or } f(c)f(b) < 0$$

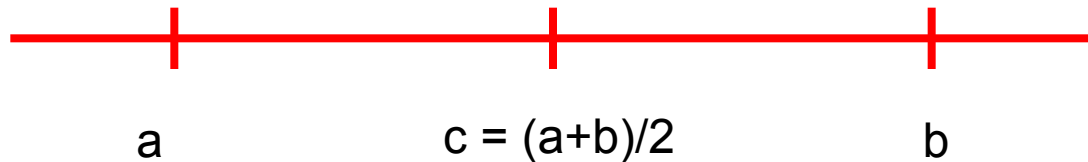
Target interval

If $f(a)f(b) < 0$

it holds that $f(a)f(c) < 0$ or $f(c)f(b) < 0$

If $f(a)f(c) < 0$, choose $[a c]$ as target interval

If $f(c)f(b) < 0$, choose $[c b]$ as target interval



Binary search

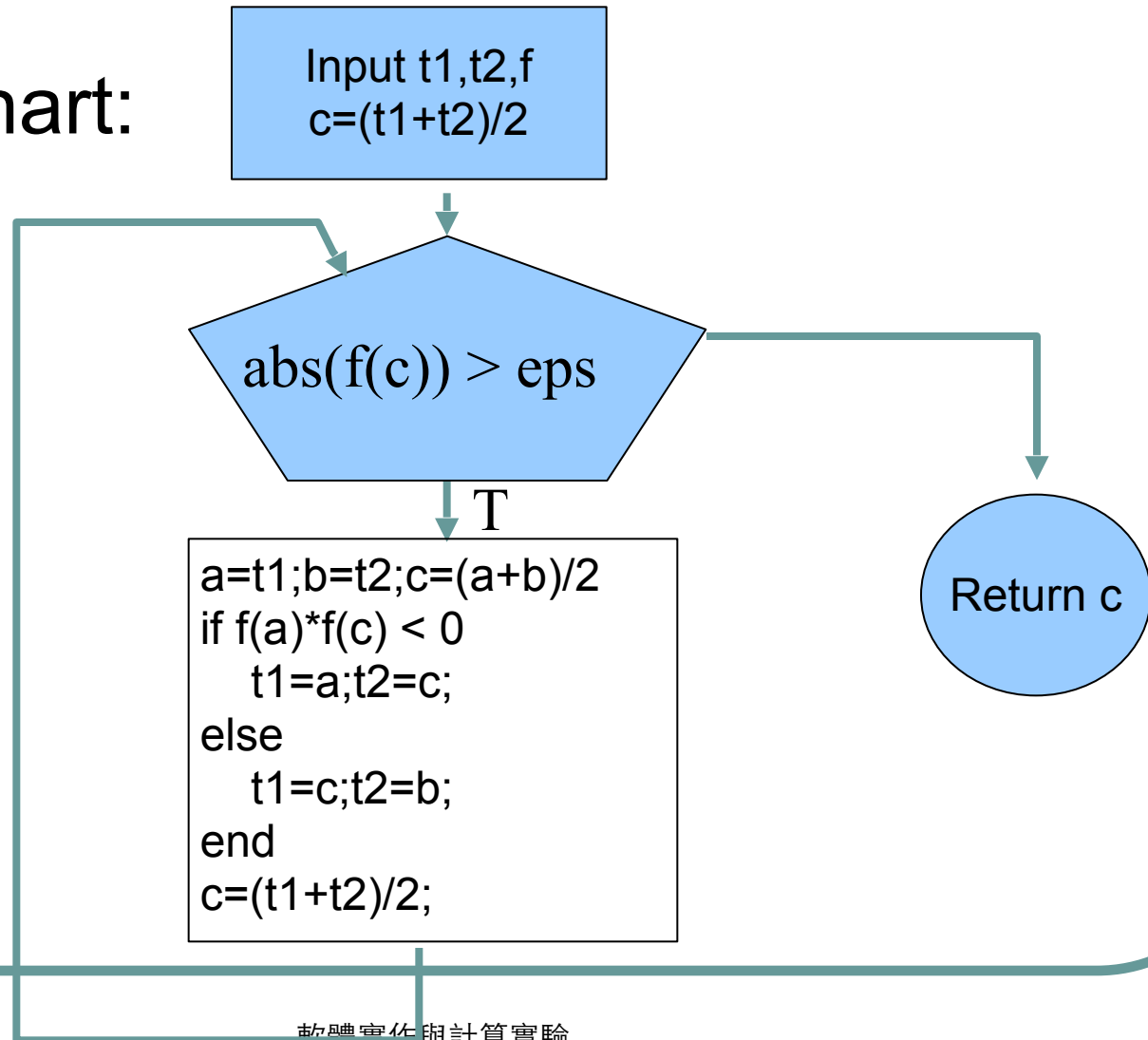
- Target interval $[t1\ t2]$
- If halting condition holds, halt
- $a \leftarrow t1, b \leftarrow t2, c \leftarrow (t1+t2)/2$
- If $f(a)f(c) < 0$, $t1 \leftarrow a, t2 \leftarrow c$ choose $[a\ c]$
- If $f(c)f(b) < 0$, $t1 \leftarrow c, t2 \leftarrow b$ choose $[c\ b]$

Halting condition

- $c = (t_1 + t_2) / 2$
- $f(c)$ is close enough to zero
- Implementation
 - $\text{abs}(f(c)) < \text{eps}$

Zero finding

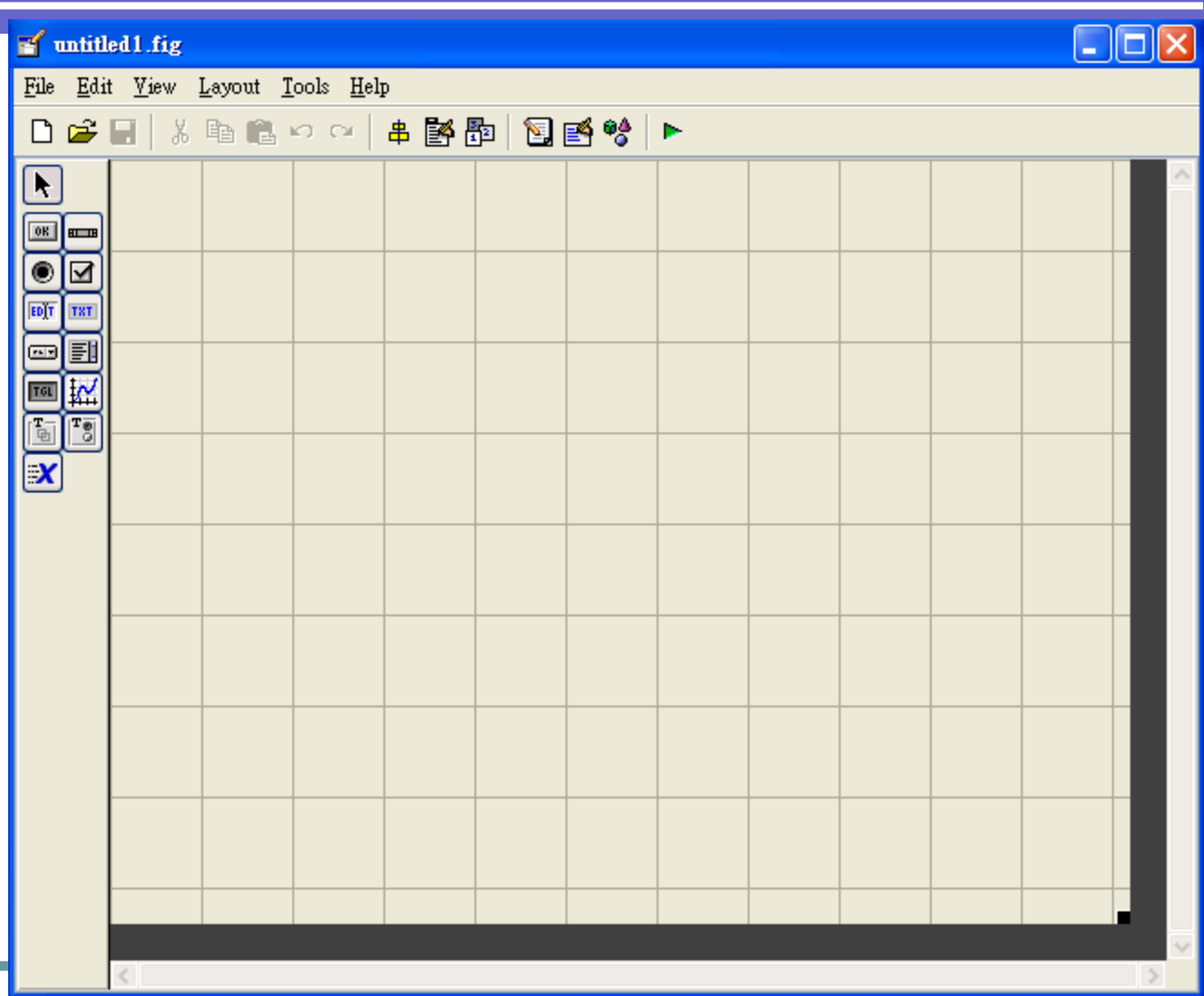
- Flow Chart:



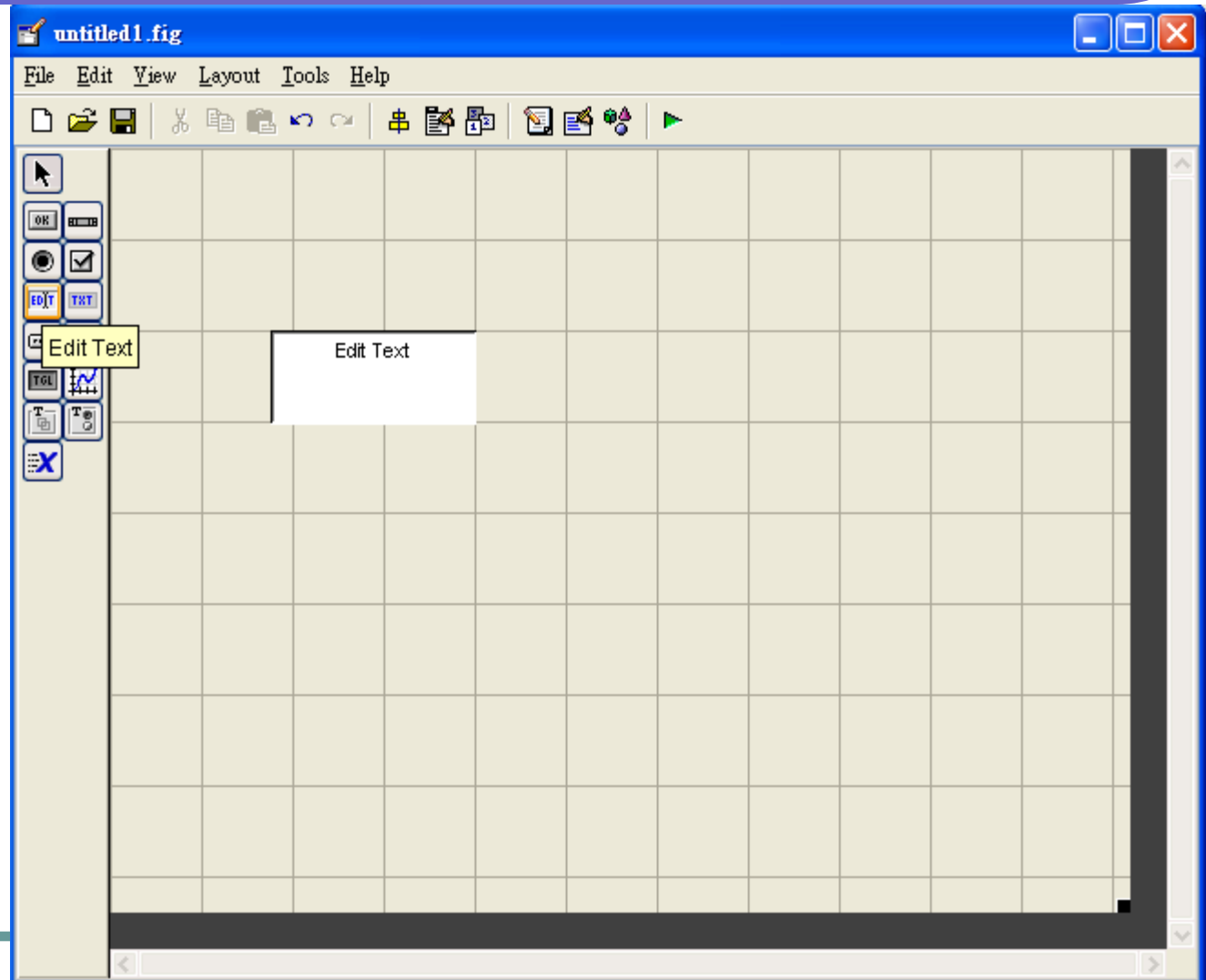
my_fzero.m

```
c=(t1+t2)/2;
while abs(f(c))>eps
    a=t1;b=t2;
    if f(a)*f(c) < 0
        t1=a;t2=c;
    else
        t1=c;t2=b;
    end
    c=(t1+t2)/2;
end
```

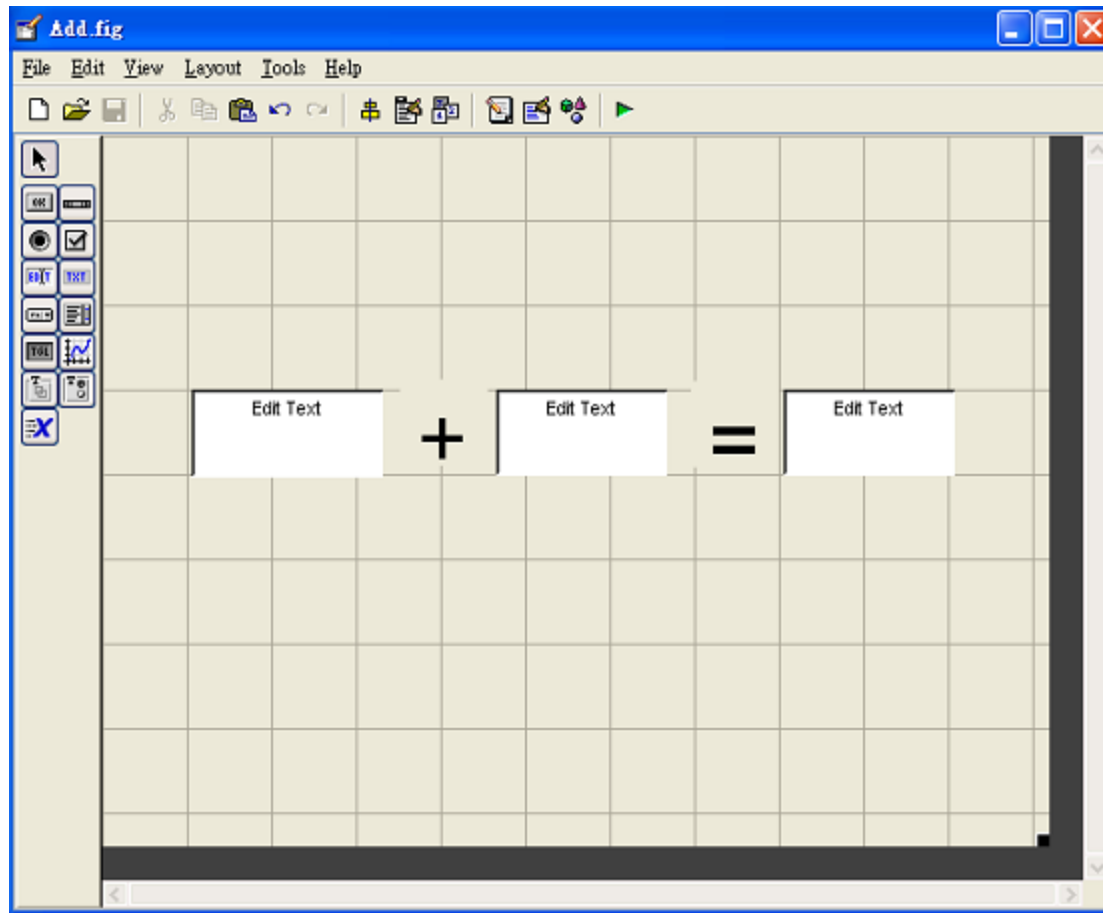
New

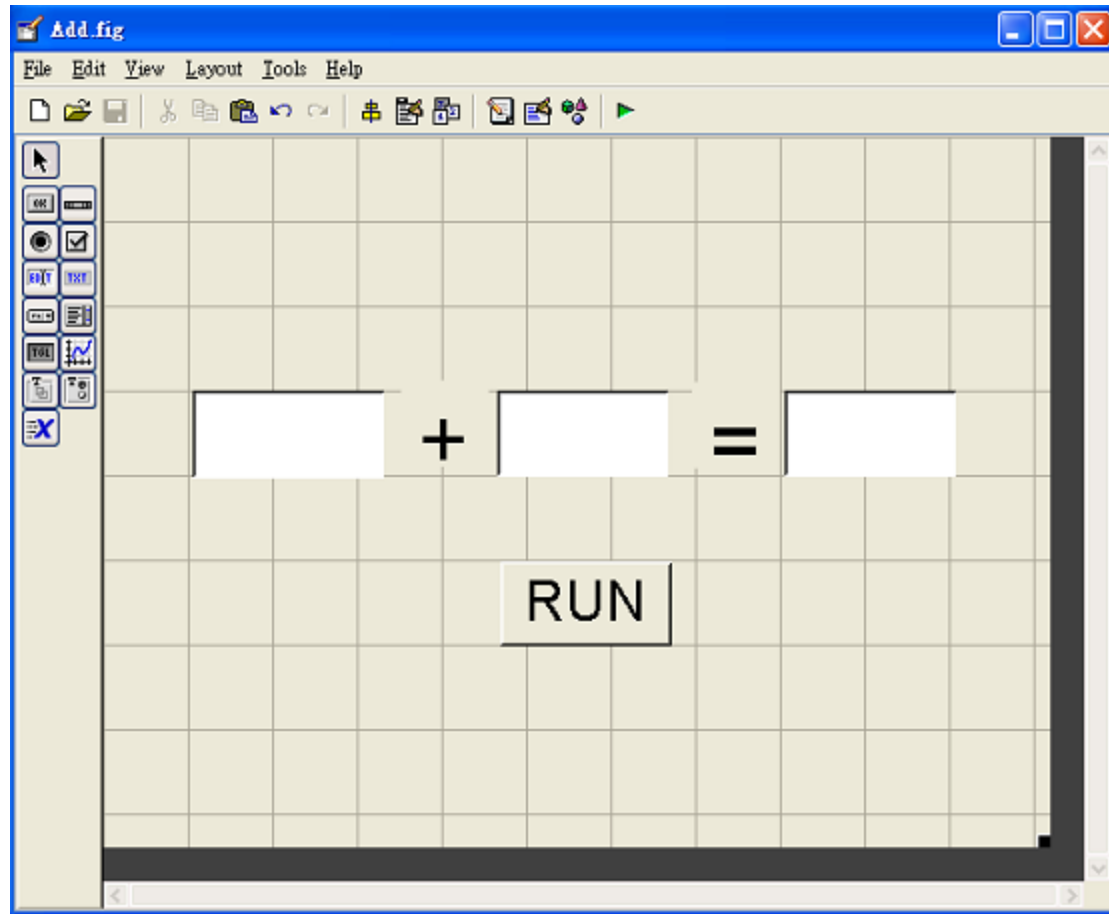


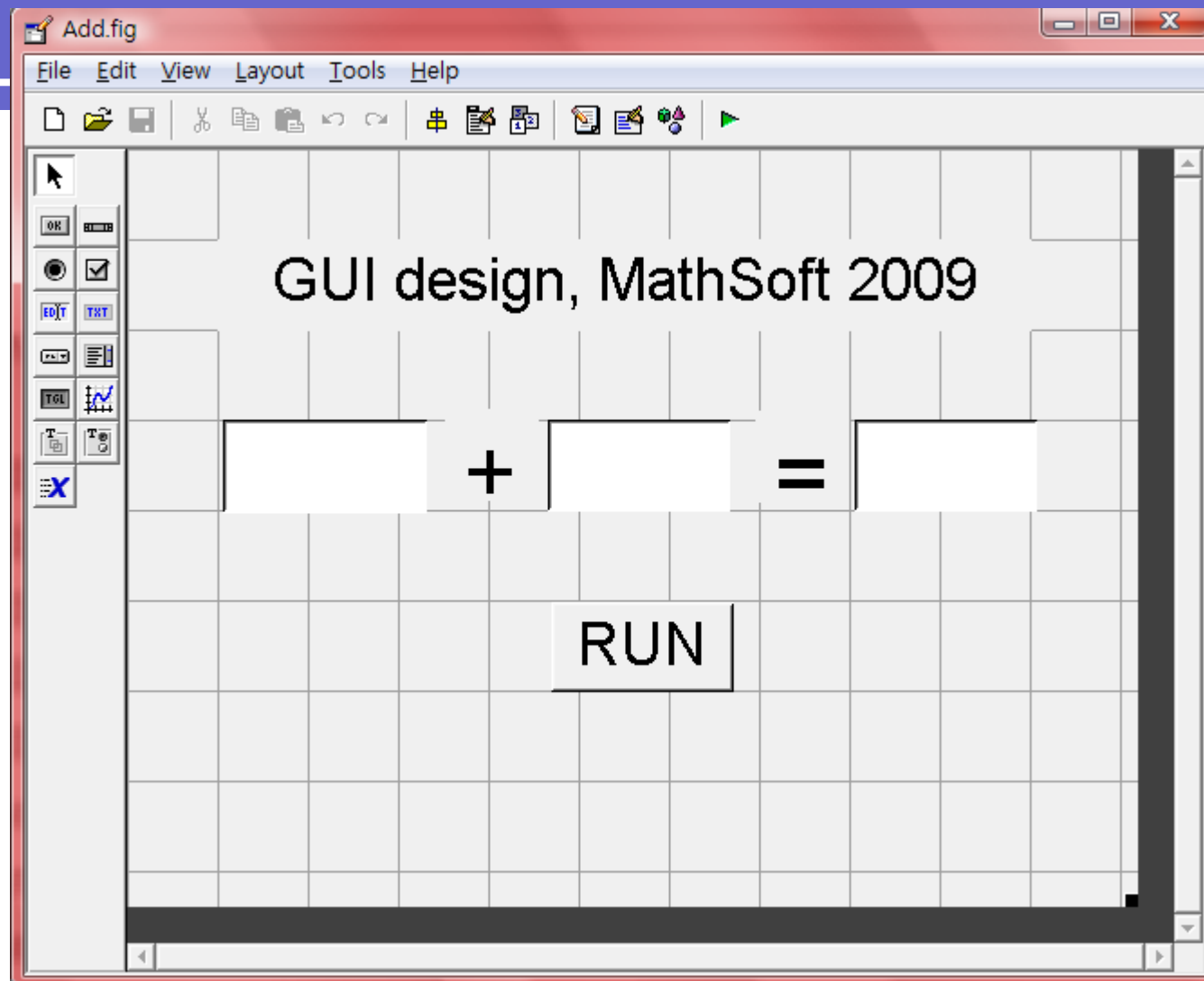
Edit Text



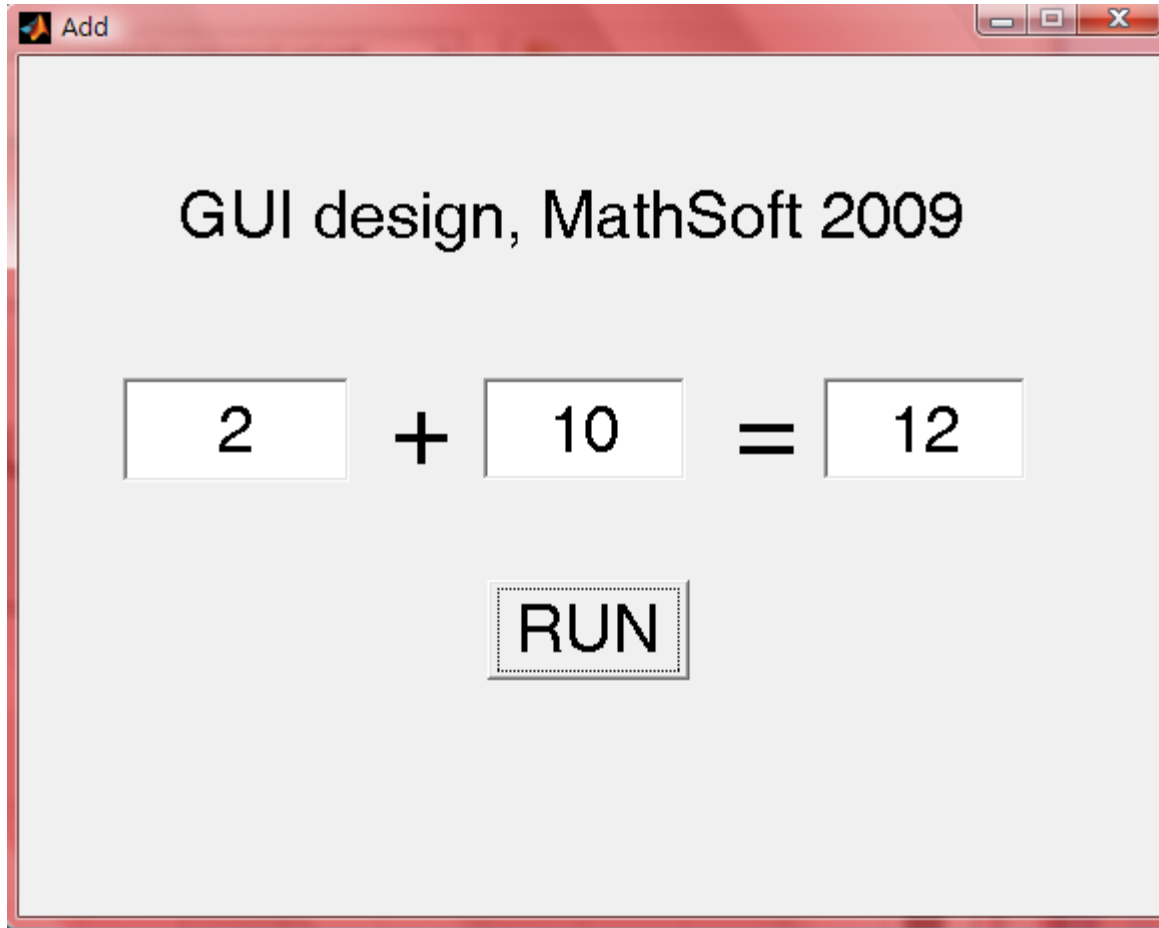
Addition







```
function pushbutton1_Callback(hObject, eventdata, handles)
% hObject    handle to pushbutton1 (see GCBO)
% eventdata  reserved - to be defined in a future version of MATLAB
% handles    structure with handles and user data (see GUIDATA)
s1=get(handles.edit1,'String');
s2=get(handles.edit2,'String');
x=str2double(s1)+str2double(s2);
s3=num2str(x);
set(handles.edit3,'String',s3);
return
```

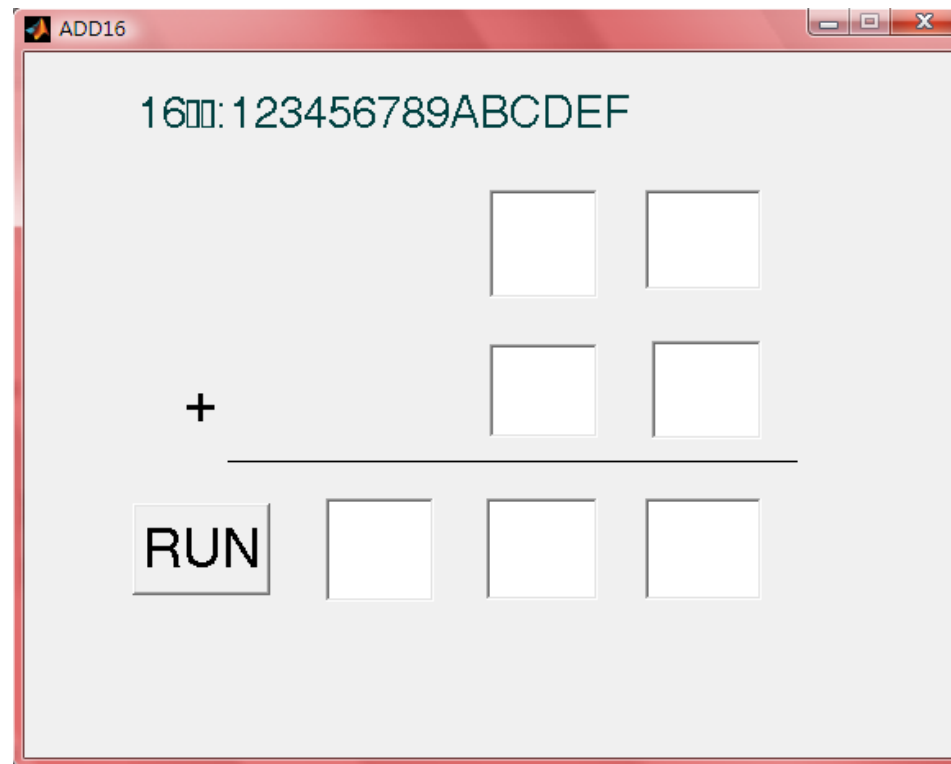
GUI tool

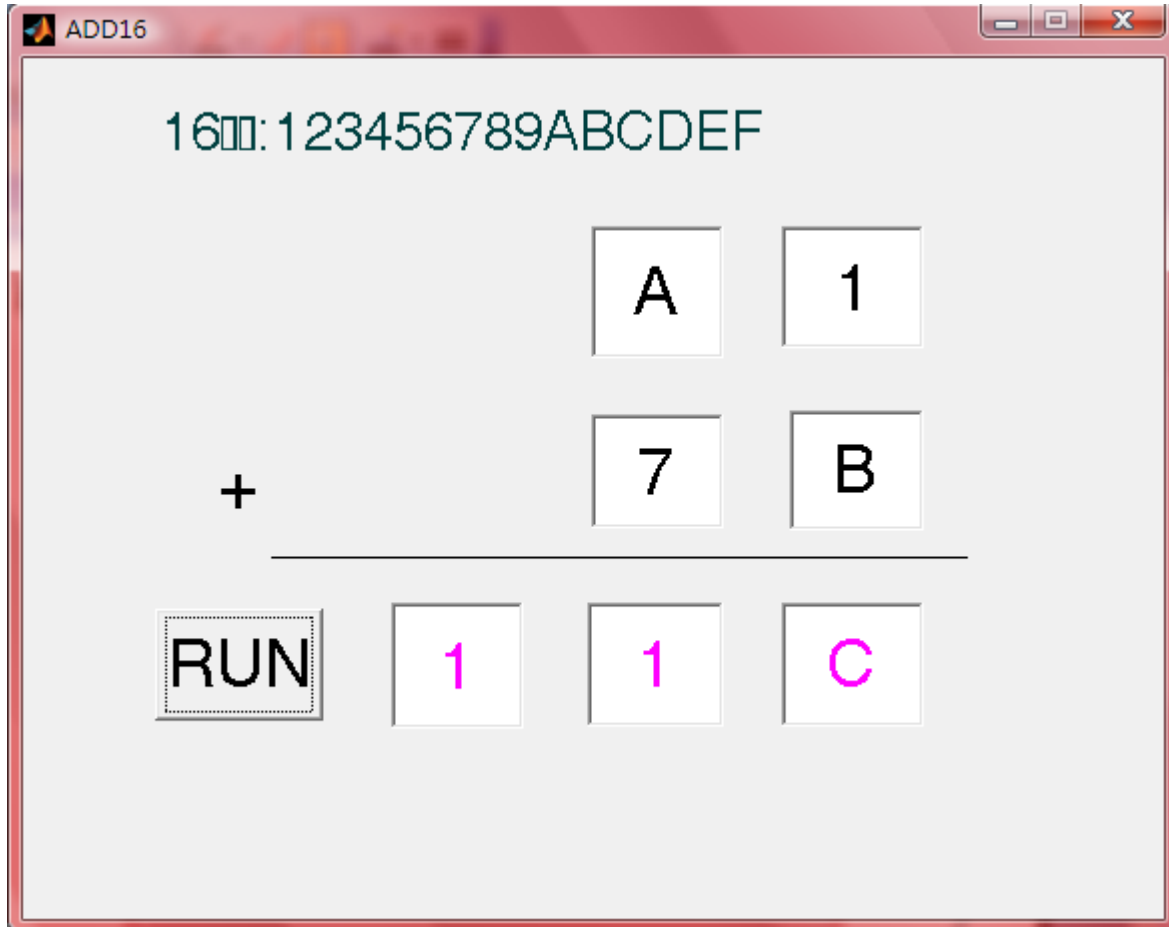
Add.fig

Add.m

ADD16

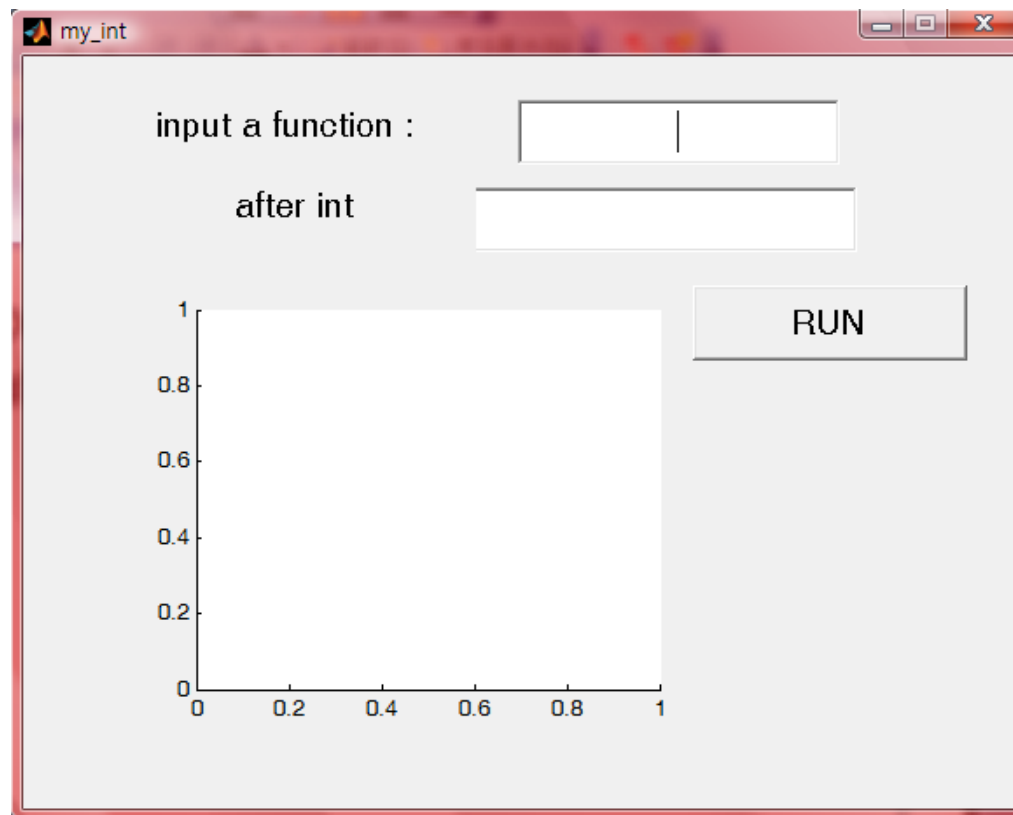
A toolbox for addition of hexadecimal numbers

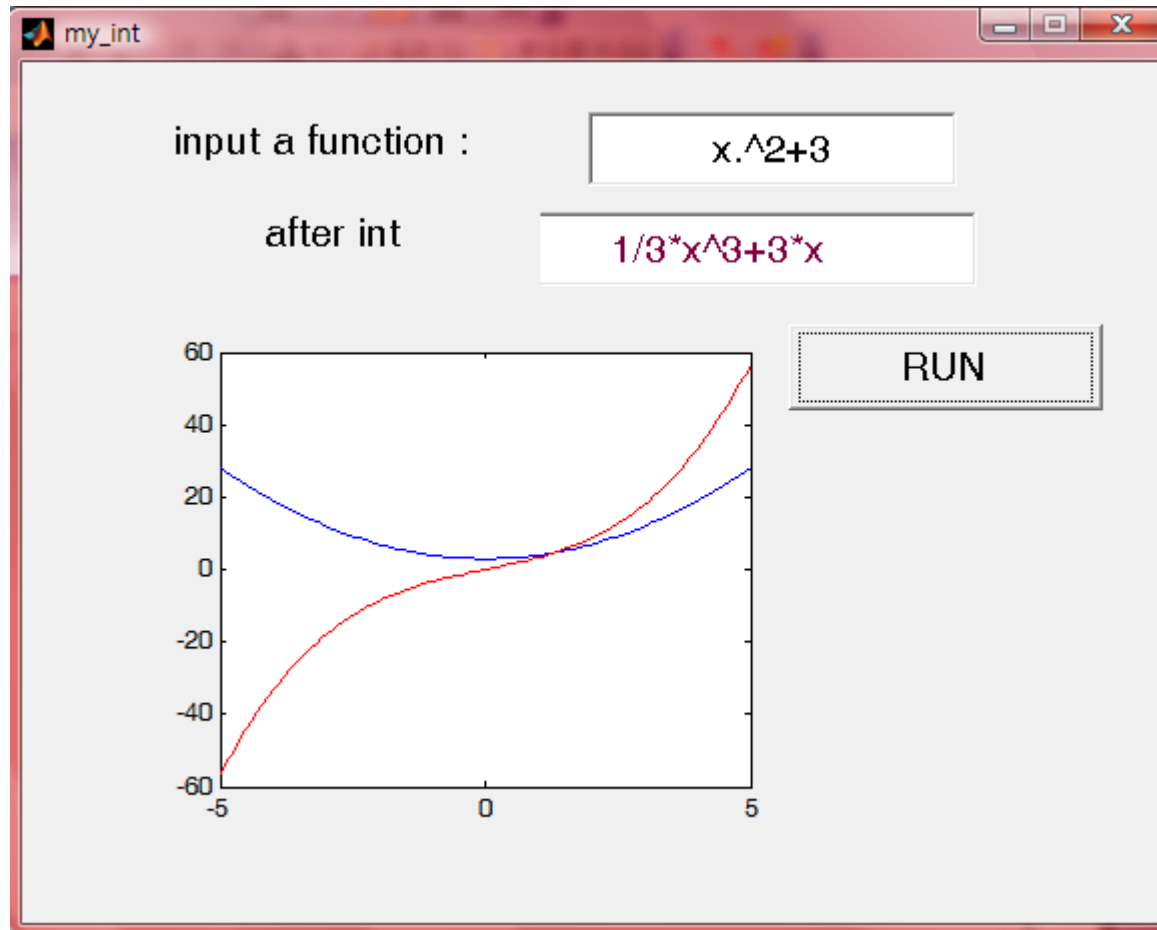






Symbolic integration

- A toolbox for demonstrating symbolic integration





Matlab Application Platform

- Web server
- Stand-alone execution
- Mobile connection
- Matlab  C++ or C  mobile app

Exercise #a#b

- Generate four distinct characters within {'0', '1', '2', '3', '4', '5', '6', '7', '8', '9'} randomly
- Draw a while-loop to realize the game of #a#b
 - Allow a player to key-in a four-digit string
 - Print n_a 'a' n_b 'b' in response to the given string
 - Halt if the guess is scored as 4'a'

- n_a denotes the number of guessed characters that appear at right position in the target
- n_b denotes the number of guessed characters that appear at wrong position in the target

Example

- Target : 6481
 - Guess : 1628 Output: 0a3b
 - Guess : 1946 Output: 0a3b
 - Guess : 6283 Output: 1a1b
 - Guess : 6481 Output: 4a0b

- Draw a flow chart to illustrate how to determine n_a for given target and guess
- Draw a flow chart to illustrate how to determine n_b for given target and guess

- Write MATLAB functions to implement flow charts for #a#b

Characters & integers

```
tt = randperm(10)-1  
cc=input('keyin:', 's');  
tt(1) == cc(1)  
tt(1) == cc(1) - '0'
```